

Lecture 6b: Phase Transition in Random Graphs

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Analytical Results

Distributions

Today we discuss about phase transition in random graphs. Recall on the *Erdős-Rényi class* $\mathcal{G}_{n,p}$ of random graphs, the probability mass function on \mathcal{G} , $P : \mathcal{G} \rightarrow [0, 1]$, is obtained by assuming that, as random variables, edges are independent from one another, and each edge occurs with probability $p \in [0, 1]$. Thus a graph $G \in \mathcal{G}$ with m vertices will have probability $P(G)$ given by

$$P(G) = p^m(1 - p)^{\binom{n}{2}}.$$

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$$P(G) = p^m (1 - p)^{\binom{n}{2} - m}.$$

Recall the expected number of q -cliques X_q is

$$\mathbb{E}[X_q] = \binom{n}{q} p^{q(q-1)/2}$$

Analytical Results

Distributions

We shall also use $\Gamma^{n,m}$ the set of all graphs on n vertices with m edges.
The set $\Gamma^{n,m}$ has cardinal

$$\binom{\binom{n}{2}}{m}.$$

In $\Gamma^{n,m}$ each graph is equally probable.

Analytical Results

Cliques

The case of 3-cliques: $\mathbb{E}[X_3] = \theta n^3 p^3$ ($\theta \sim \frac{1}{6}$).

The case of 4-cliques: $\mathbb{E}[X_4] = \theta n^4 p^6$ ($\theta \sim \frac{1}{24}$).

The first problem we consider is the size of the largest clique of a random graph.

Note, finding the size of the largest clique (called *the clique number*) is a NP-hard problem.

Analytical Results

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Idea: Analyze p so that $\mathbb{E}[X_q] \approx 1$.

- For $p > \frac{1}{n}$ and large n we expect that graphs will have a 3-clique;
- For $p > \frac{1}{n^{2/3}}$ and large n , we expect that graphs will have a 4-clique;

Analytical Results

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- For $p > \frac{1}{n}$ and large n we expect that graphs will have a 3-clique;
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Question: How sharp are these thresholds?

Analytical Results

3-Cliques

Theorem

Let $p = p(n)$ be the edge probability in $\mathcal{G}_{n,p}$.

- 1 If $p \gg \frac{1}{n}$ (i.e. $\lim_{n \rightarrow \infty} np = \infty$) then
 $\lim_{n \rightarrow \infty} \text{Prob}[G \in \mathcal{G}_{n,p} \text{ has a 3-clique}] \rightarrow 1.$
- 2 If $p \ll \frac{1}{n}$ (i.e. $\lim_{n \rightarrow \infty} np = 0$) then
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Let $m = m(n)$ be the number of edges in $\Gamma^{n,m}$.

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Analytical Results

4-Cliques

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Let $p = p(n)$ be the edge probability in $\mathcal{G}_{n,p}$.

- 1 If $p \gg \frac{1}{n^{2/3}}$ (i.e. $\lim_{n \rightarrow \infty} n^{2/3}p = \infty$) then $\lim_{n \rightarrow \infty} \text{Prob}[G \in \mathcal{G}_{n,p} \text{ has a 4-clique}] \rightarrow 1$.
- 2 If $p \ll \frac{1}{n^{2/3}}$ (i.e. $\lim_{n \rightarrow \infty} n^{2/3}p = 0$) then $\lim_{n \rightarrow \infty} \text{Prob}[G \in \mathcal{G}_{n,p} \text{ has a 4-clique}] \rightarrow 0$.

Analytical Results

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Let $p = p(n)$ be the edge probability in $\mathcal{G}_{n,p}$.

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Theorem

Let $m = m(n)$ be the number of edges in $\Gamma^{n,m}$.

- 1 If $m \gg n^{4/3}$ (i.e. $\lim_{n \rightarrow \infty} \frac{m}{n^{4/3}} = \infty$) then $\lim_{n \rightarrow \infty} \text{Prob}[G \in \Gamma^{n,m} \text{ has a 4-clique}] \rightarrow 1$.
- 2 If $m \ll n^{4/3}$ (i.e. $\lim_{n \rightarrow \infty} \frac{m}{n^{4/3}} = 0$) then $\lim_{n \rightarrow \infty} \text{Prob}[G \in \Gamma^{n,m} \text{ has a 4-clique}] \rightarrow 0$.

Analytical Results

q -Cliques

Theorem

Let $p = p(n)$ be the edge probability in $\mathcal{G}_{n,p}$. Let $q \geq 3$ be an integer.

- 1 If $p \gg \frac{1}{n^{2/(q-1)}}$ (i.e. $\lim_{n \rightarrow \infty} n^{2/(q-1)} p = \infty$) then $\lim_{n \rightarrow \infty} \text{Prob}[G \in \mathcal{G}_{n,p} \text{ has a } q\text{-clique}] \rightarrow 1$.
- 2 If $p \ll \frac{1}{n^{2/(q-1)}}$ (i.e. $\lim_{n \rightarrow \infty} n^{2/(q-1)} p = 0$) then $\lim_{n \rightarrow \infty} \text{Prob}[G \in \mathcal{G}_{n,p} \text{ has a } q\text{-clique}] \rightarrow 0$.

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Theorem

Let $m = m(n)$ be the number of edges in $\Gamma^{n,m}$. Let $q \geq 3$ be and integer.

- 1 If $m \gg n^{2(q-2)/(q-1)}$ (i.e. $\lim_{n \rightarrow \infty} \frac{m}{n^{2(q-2)/(q-1)}} = \infty$) then
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- 2 If $m \ll n^{2(q-2)/(q-1)}$ (i.e. $\lim_{n \rightarrow \infty} \frac{m}{n^{2(q-2)/(q-1)}} = 0$) then
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Analytical Results

Markov and Chebyshev Inequalities

We want to control probabilities of the random event $X_3(G) > 0$. Two important tools:

- 1 (Markov's Inequality) Assume X is a non-negative random variable. Then $Prob[X \geq t] \leq \frac{\mathbb{E}[X]}{t}$.
- 2 (Chebyshev's Inequality) For any random variable X , $Prob[|X - E[X]| \geq t] \leq \frac{Var[X]}{t^2}$.

where $\mathbb{E}[X]$ is the mean of X , and $Var[X] = \mathbb{E}[X^2] - |\mathbb{E}[X]|^2$ is the variance of X .

Analytical Results

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where $\mathbb{E}[X]$ is the mean of X , and $\text{Var}[X] = \mathbb{E}[X^2] - |\mathbb{E}[X]|^2$ is the variance of X . *Quick Proof:*

$$\text{Prob}[X \geq t] = \int_t^\infty p_X(x) dx \leq \frac{1}{t} \int_t^\infty x p_X(x) dx \leq \frac{\mathbb{E}[X]}{t}.$$

$$\text{Prob}[|X - \mathbb{E}[X]| \geq t] = P[|X - \mathbb{E}[X]|^2 \geq t^2] \leq \frac{\mathbb{E}[|X - \mathbb{E}[X]|^2]}{t^2} = \frac{\text{Var}[X]}{t^2}.$$

Analytical Results

Proofs for the 3-clique case

For small probability: We shall use Markov's inequality to show

$Prob[X_3 > 0] \rightarrow 0$ when $p \ll \frac{1}{n}$:

$$Prob[X_3 > 0] = Prob[X_3 \geq 1] \leq \frac{E[X_3]}{1} = \frac{n(n-1)(n-2)}{6} p^3 = \theta n^3 p^3 \rightarrow 0.$$

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For large probability: Since $\mathbb{E}[X_3] \rightarrow \infty$ it follows that $Prob[X_3 > 0] > 0$.

We need to show that $Prob[X_3 = 0] \rightarrow 0$. By Chebyshev's inequality:

$$Prob[X_3 = 0] \leq Prob[|X_3 - \mathbb{E}[X_3]| \geq \mathbb{E}[X_3]] \leq \frac{Var[X_3]}{|\mathbb{E}[X_3]|^2}$$

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Need the variance of $X_3 = \sum_{(i,j,k) \in S_3} 1_{i,j,k}$,

$$X_3^2 = \sum_{(i,j,k) \in S_3} \sum_{(i',j',k') \in S_3} 1_{i,j,k} 1_{i',j',k'}.$$

Analytical Results

Proofs for the 3-clique case

$$\begin{aligned}
 X_3^2 = & \sum_{(i,j,k) \in \mathcal{S}_3(n)} 1_{i,j,k} + \sum_{(i,j,k) \in \mathcal{S}_3(n)} \sum_{l \in \mathcal{S}_1(n-3)} (1_{i,j,k} 1_{i,j,l} + 1_{i,j,k} 1_{j,k,l} + 1_{i,j,k} 1_{k,i,l}) + \\
 & + \sum_{(i,j,k) \in \mathcal{S}_3(n)} \sum_{u,v \in \mathcal{S}_2(n-3)} (1_{i,j,k} 1_{i,u,v} + 1_{i,j,k} 1_{j,u,v} + 1_{i,j,k} 1_{k,u,v}) + \\
 & + \sum_{(i,j,k) \in \mathcal{S}_3(n)} \sum_{(i',j',k') \in \mathcal{S}_3(n-3)} 1_{i,j,k} 1_{i',j',k'}
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 & + \sum_{(i,j,k) \in \mathcal{S}_3(n)} \sum_{(i',j',k') \in \mathcal{S}_3(n-3)} 1_{i,j,k} 1_{i',j',k'}
 \end{aligned}$$

$$\mathbb{E}[X_3^2] = |\mathcal{S}_3| p^3 + 3|\mathcal{S}_3|(n-3)p^5 + 3|\mathcal{S}_3| \binom{n-3}{2} p^6 + |\mathcal{S}_3| \binom{n-3}{3} p^6.$$

Thus

$$\text{Var}[X_3] = \mathbb{E}[X_3^2] - |\mathbb{E}[X_3]|^2 = \dots = \theta(n^3 p^3 + n^4 p^5 + n^5 p^6).$$

Analytical Results

Proofs for the 3-clique case

and:

$$\text{Prob}[X_3 = 0] \leq \frac{\theta(n^3 p^3 + n^4 p^5 + n^5 p^6)}{\theta(n^6 p^6)} = \frac{1}{(np)^3} + \frac{1}{n} \rightarrow 0.$$

Analytical Results

Proofs for the 3-clique case

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Similar proofs for the other cases (4-cliques and q -cliques).

Analytical Results

Behavior at the threshold

In general we obtain a "coarse threshold". Recall a Poisson random variable X with parameter λ has p.m.f. $Prob[X = k] = e^{-\lambda} \frac{\lambda^k}{k!}$.

Theorem

In $\mathcal{G}_{n,p}$,

- ① For $p = \frac{c}{n}$, X_3 is asymptotically Poisson with parameter $\lambda = c^3/6$.
- ② For $p = \frac{c}{n^{2/3}}$, X_4 is asymptotically Poisson with parameter $\lambda = c^6/24$.

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Theorem

In $\Gamma^{n,m}$,

- ① For $m = cn$, X_3 is asymptotically Poisson with parameter $\lambda = 4c^3/3$.
- ② For $m = cn^{4/3}$, X_4 is asymptotically Poisson with parameter $\lambda = 8c^6/3$.

Analytical Results

Connected Components

$\mathcal{G}_{n,p}$ class of random graphs has a remarkable property in regards to the largest connected component. We shall express the result in the class $\Gamma^{n,m}$.

Analytical Results

Connected Components

Theorem

- ① Let $m = m(n)$ satisfies $m \ll \frac{1}{2}n \log(n)$. Then

$$\lim_{n \rightarrow \infty} \text{Prob}[G \in \Gamma^{n,m} \text{ is connected}] = 0$$

- ② Let $m = m(n)$ satisfies $m \gg \frac{1}{2}n \log(n)$. Then

$$\lim_{n \rightarrow \infty} \text{Prob}[G \in \Gamma^{n,m} \text{ is connected}] = 1$$

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- ③ Assume $m = \frac{1}{2}n \log(n) + tn + o(n)$, where $o(n) \ll n$. Then

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In this case $\frac{1}{2}n \log(n)$ is known as a *strong threshold*. 

Numerical Results

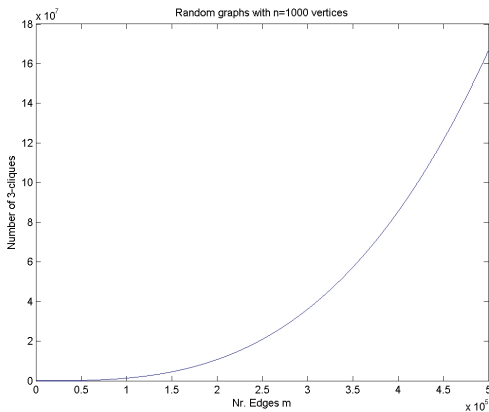
3-cliques & Connectivity results

Results for $n = 1000$ vertices.

- 1 3-cliques. Recall $\mathbb{E}[X_3] \sim m^3$
- 2 Connectivity. Recall the connectivity threshold is $\frac{1}{2}n \log(n) = 3454$.

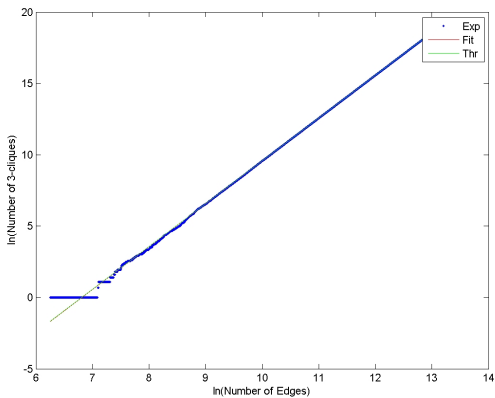
Numerical Results

3-cliques



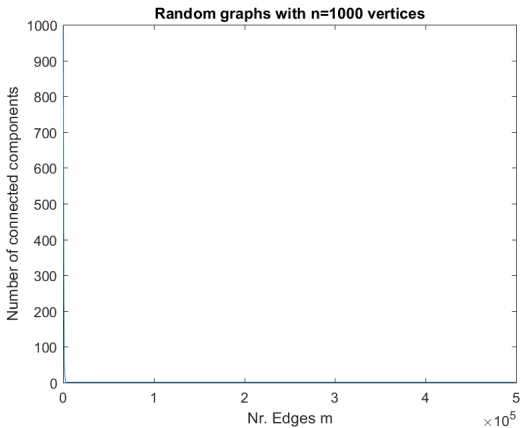
Numerical Results

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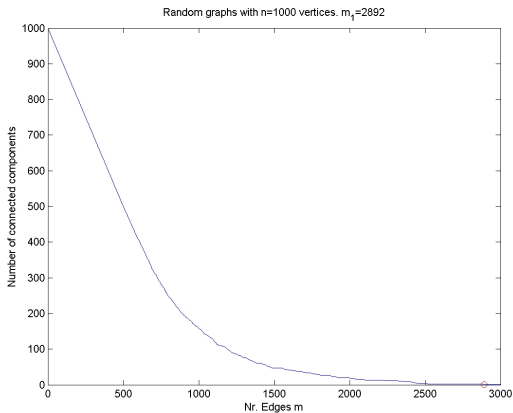
Numerical Results

Connectivity



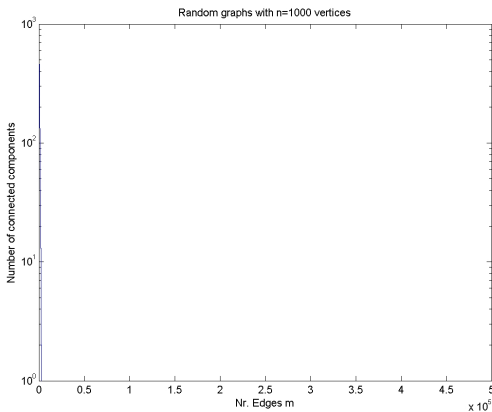
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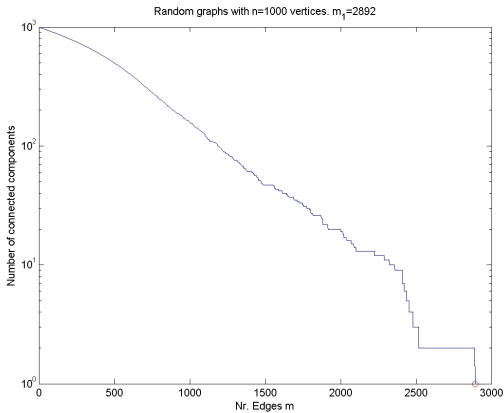
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Numerical Results

Connectivity



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