# Lecture 3: Random Graphs

#### Radu Balan

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# The Erdös-Rényi class $\mathcal{G}_{n,p}$

Today we discuss about random graphs. The *Erdös-Rényi class*  $\mathcal{G}_{n,p}$  of random graphs is defined as follows.

Algorithmics

# The Erdös-Rényi class $\mathcal{G}_{n,p}$

Today we discuss about random graphs. The *Erdös-Rényi class*  $\mathcal{G}_{n,p}$  of random graphs is defined as follows. Let  $\mathcal{V}$  denote the set of *n* vertices,  $\mathcal{V} = \{1, 2, \dots, n\}$ , and let  $\mathcal{G}$  denote the set of all graphs with vertices  $\mathcal{V}$ . There are exactly  $2 \begin{pmatrix} n \\ 2 \end{pmatrix}$  such graphs. The probability mass function on  $\mathcal{G}$ ,  $P : \mathcal{G} \to [0, 1]$ , is obtained by assuming that, as random variables, edges are independent from one another, and each edge occurs with probability P(G) given by

$$P(G) = p^{m}(1-p)^{\binom{n}{2}-m}$$

(explain why)

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## The Erdös-Rényi class $\mathcal{G}_{n,p}$ Probability space

Formally,  $\mathcal{G}_{n,p}$  stands for the the probability space  $(\mathcal{G}, P)$  composed of the set  $\mathcal{G}$  of all graphs with *n* vertices, and the probability mass function *P* defined above.

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A reformulation of *P*: Let  $G = (\mathcal{V}, \mathcal{E})$  be a graph with *n* vertices and *m* edges and let *A* be its adjacency matrix. Then:

$$P(G) = \prod_{(i,j)\in\mathcal{E}} Prob((i,j) \text{ is an edge}) \prod_{(i,j)\notin\mathcal{E}} Prob((i,j) \text{ is not an edge}) =$$
$$= \prod_{1\leq i< j\leq n} p^{A_{i,j}} (1-p)^{1-A_{i,j}}$$

where the product is over all ordered pairs (i, j) with  $1 \le i < j \le n$ . Note:

$$|\{(i,j), 1 \le i < j \le n\}| = \binom{n}{2} \& |\{(i,j) \in \mathcal{E}\}| = |\mathcal{E}| = m = \sum_{1 \le i < j \le n} A_{i,j}.$$

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## The Erdös-Rényi class $\mathcal{G}_{n,p}$ Computations in $\mathcal{G}_{n,p}$

## How to compute the expected number of edges of a graph in $\mathcal{G}_{n,p}$ ?

Image: A matrix and a matrix

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## The Erdös-Rényi class $\mathcal{G}_{n,p}$ Computations in $\mathcal{G}_{n,p}$

How to compute the expected number of edges of a graph in  $\mathcal{G}_{n,p}$ ? Let  $X_2 : \mathcal{G}_{n,p} \to \{0, 1, \cdots, \binom{n}{2}\}$  be the random variable of *number of edges of a graph G*.

$$X_2 = \sum_{1 \le i < j \le n} 1_{(i,j)}$$
,  $1_{(i,j)}(G) = \begin{cases} 1 & \text{if } (i,j) \text{ is edge in } G \\ 0 & \text{if } otherwise \end{cases}$ 

Use linearity and the fact that  $\mathbb{E}[1_{(i,j)}] = Prob((i,j) \in \mathcal{E}) = p$  to obtain:

$$\mathbb{E}[\text{Number of Edges}] = \left( egin{array}{c} n \\ 2 \end{array} 
ight) p = rac{n(n-1)}{2}p$$

## The Erdös-Rényi class $\mathcal{G}_{n,p}$ MLE of p

Given a realization G of a graph with n vertices and m edges, how to estimate the most likely p that explains the graph. Concept: The Maximum Likelihood Estimator (MLE). In statistics: The MLE of a parameter  $\theta$  given an observation x of a random variable  $X \sim p_X(x; \theta)$  is the value  $\theta$  that maximizes the probability  $P_X(x; \theta)$ :

$$\theta_{MLE} = \operatorname{argmax}_{\theta} P_X(x; \theta).$$

In our case: our observation G has m edges. We know

$$P(G; p) = p^{m}(1-p) \binom{n}{2}^{-m}$$

## The Erdös-Rényi class $\mathcal{G}_{n,p}$ MLE of p

#### Lemma

Given a random graph with n vertices and m edges, the MLE estimator of p is

$$p_{MLE} = \frac{m}{\binom{n}{2}} = \frac{2m}{n(n-1)}.$$

## The Erdös-Rényi class $\mathcal{G}_{n,p}$ MLE of p

#### Lemma

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Why

Note 
$$log(P(G; p)) = mlog(p) + \left(\binom{n}{2} - m\right)log(1-p)$$
 and solve for  $p$ :  

$$\frac{dlog(P)}{dp} = \frac{m}{p} - \frac{\binom{n}{2} - m}{1-p} = 0.$$
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## The Erdös-Rényi class $\mathcal{G}_{n,p}$ Method of Moments Estimator for p

An alternative parameter estimation method is the moment matching method. Given a likelihood function for observed data  $p(x; \theta)$  and a sequence of observations  $(x_1, x_2, \dots, x_N)$ , the moment matching method computes the parameters  $\theta \in \mathbb{R}^d$  by solving the system of equations:

$$\mathbb{E}[X] = \frac{1}{N} \sum_{t=1}^{N} x_t \cdots \mathbb{E}[X^d] = \frac{1}{N} \sum_{t=1}^{N} x_t^d$$

(or unbiased estimates of the moments). In particular, for the Erdös-Rényi class, we match the first moment with the observation:  $\frac{n(n-1)}{2}p = m$ . Hence

$$p_{MM}=\frac{2m}{n(n-1)},$$

same as the MLE estimator.

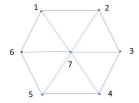
## Cliques q-cliques

### Definition

Given a graph  $G = (\mathcal{V}, \mathcal{E})$ , a subset of q vertices  $S \subset \mathcal{V}$  is called a q-clique if the subgraph  $(S, \mathcal{E}|_S)$  is complete.

In other words, S is a q-clique if for every  $i \neq j \in S$ ,  $(i, j) \in \mathcal{E}$  (or  $(j, i) \in \mathcal{E}$ ), that is, (i, j) is an edge in G.

• Each edge is a 2-clique.

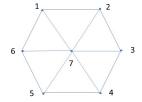


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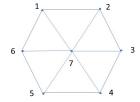
- Each edge is a 2-clique.
- {1,2,7} is a 3-clique. And so are {2,3,7}, {3,4,7}, {4,5,7}, {5,6,7}, {1,6,7}

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- There is no k-clique, with  $k \ge 4$ .

## The Erdös-Rényi class $\mathcal{G}_{n,p}$ Computations in $\mathcal{G}_{n,p}$ : *q*-cliques

How to compute the expected number of *q*-cliques?

Image: A matrix and a matrix

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## The Erdös-Rényi class $\mathcal{G}_{n,p}$ Computations in $\mathcal{G}_{n,p}$ : *q*-cliques

How to compute the expected number of *q*-cliques?

For k = 2 we computed earlier the number of edges, which is also the number of 2-cliques.

We shall compute now the number of 3-cliques: triangles, or 3-cycles.

Let  $X_3 : \mathcal{G}_{n,p} \to \mathbb{N}$  be the random variable of number of 3-cliques. Note

the maximum number of 3-cliques is  $\begin{pmatrix} n \\ 3 \end{pmatrix}$ .

Let  $S_3$  denote the set of all distinct 3-cliques of the complete graph with n vertices,  $S_3 = \{(i, j, k) , 1 \le i < j < k \le n\}$ . Let

$$1_{(i,j,k)}(G) = \begin{cases} 1 & if \quad (i,j,k) \text{ is a } 3-clique \text{ in } G\\ 0 & if \quad otherwise \end{cases}$$

## The Erdös-Rényi class $\mathcal{G}_{n,p}$ Expectation of the number of 3-cliques

Note: 
$$X_3 = \sum_{(i,j,k) \in S_3} 1_{(i,j,k)}$$
. Thus  
 $\mathbb{E}[X_3] = \sum_{(i,j,k) \in S_3} \mathbb{E}[1_{(i,j,k)}] = \sum_{(i,j,k) \in S_3} Prob((i,j,k) \text{ is a clique}).$ 

Since  $Prob((i, j, k) \text{ is a clique}) = p^3$  we obtain:

$$\mathbb{E}[\text{Number of } 3-\text{cliques}] = \binom{n}{3}p^3 = \frac{n(n-1)(n-2)}{6}p^3.$$

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## The Erdös-Rényi class $\mathcal{G}_{n,p}$ Number of 3 cliques

Assume we observe a graph G with n vertices and m edges. What would be the expected number  $N_3$  of 3-cliques?

$$\mathbb{E}[X_3|X_2 = m] = \frac{1}{L} \sum_{k=1}^{L} X_3(G_k)$$

where *L* denotes the numbe of graphs with *m* edges and *n* vertices, and  $G_1, \dots, G_L$  is an enumeration of these graphs.

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where *L* denotes the numbe of graphs with *m* edges and *n* vertices, and  $G_1, \dots, G_L$  is an enumeration of these graphs. We approximate:

$$\mathbb{E}[X_3|X_2=m] \approx \mathbb{E}[X_3; p = p_{MLE}(m)]$$

and obtain:

$$E[X_3|X_2=m] \approx \frac{4(n-2)}{3n^2(n-1)^2}m^3$$

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## The Erdös-Rényi class $\mathcal{G}_{n,p}$ Expectation of the number of *q*-cliques

Let  $X_q : \mathcal{G}_{n,p} \to \mathbb{N}$  be the random variable of number of *q*-cliques. Note the maximum number of *q*-cliques is  $\binom{n}{q}$ . Let  $S_q$  denote the set of all distinct *q*-cliques of the complete graph with *n* vertices,  $S_q = \{(i_1, i_2, \cdots, i_q), 1 \leq i_1 < i_2 < \cdots < i_q \leq n\}$ . Note  $|S_q| = \binom{n}{q}$ . Let

$$1_{(i_1,i_2,\cdots,i_q)}(G) = \begin{cases} 1 & if \quad (i_1,i_2,\cdots,i_q) \text{ is a } q-clique \text{ in } G \\ 0 & if \quad otherwise \end{cases}$$

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## The Erdös-Rényi class $\mathcal{G}_{n,p}$ Expectation of the number of *q*-cliques

Expectation of the number of q-cliques

Since 
$$X_q = \sum_{(i_1, \dots, i_q) \in S_q} 1_{i_1, \dots, i_q}$$
 and  
 $Prob((i_1, \dots, i_q) \text{ is a clique}) = p^{\begin{pmatrix} q \\ 2 \end{pmatrix}}$  we obtain:

$$\mathbb{E}[\textit{Number of } q-\textit{cliques}] = \left(egin{array}{c} n \ q \end{array}
ight) p^{q(q-1)/2}.$$

## The Erdös-Rényi class $\mathcal{G}_{n,p}$ Expectation of the number of *q*-cliques

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$$\mathbb{E}[\text{Number of } q-\text{cliques}]=\left(egin{array}{c}n\\q\end{array}
ight)p^{q(q-1)/2}.$$

Using a similar argument as before, if G has m edges, then

$$\mathbb{E}[X_q|X_2=m]\approx \binom{n}{q}\left(\frac{2m}{n(n-1)}\right)^{q(q-1)/2}$$

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## Computation of Number of Cliques An Iterative Algorithm

We discuss two algorithms to compute  $X_q$ : iterative, and adjacency matrix based algorithm.

*Framework*: we are given a sequence  $(G_t)_{t\geq 0}$  of graphs on n vertices, where  $G_{t+1}$  is obtained from  $G_t$  by adding one additional edge:  $G_t = (\mathcal{V}, \mathcal{E}_t), \ \emptyset = \mathcal{E}_0 \subset \mathcal{E}_1 \subset \cdots$  and  $|\mathcal{E}_t| = t$ . **Iterative Algorithm**: Assume we know  $X_q(G_t)$ , the number of q-cliques of graph  $G_t$ . Then  $X_q(G_{t+1}) = X_q(G_t) + D_q(e; G_t)$  where  $D_q(e; G_t)$  denotes the number of q-cliques in  $G_{t+1}$  formed by the additional edge  $e \in \mathcal{E}_{t+1} \setminus \mathcal{E}_t$ .

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## Computation of Number of Cliques An Analytic Formula

Laplace Matrix  $\Delta = D - A$  contains all connectivity information. *Idea*: Note the (i, j) element of  $A^2$  is

$$(A^2)_{i,j} = \sum_{k=1}^n A_{i,k} A_{k,j} = |\{k : i \sim k \sim j\}|.$$

This means  $(A^2)_{i,j}$  is the number of paths of length 2 that connect *i* to *j*. *Remark*: The diagonal elements of  $A(A^2 - D)$  represent twice the number of 3-cycles (= 3-cliques) that contain that particular vertex. *Conclusion*:

$$X_3 = \frac{1}{6} trace \{A(A^2 - D)\} = \frac{1}{6} trace(A^3).$$

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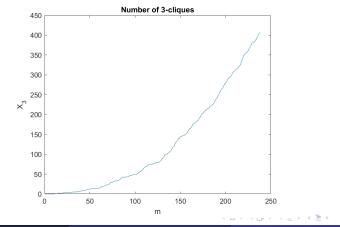
*Exercise*: Generalize this formula for  $X_4$ .

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#### Numerical results Graph of $X_3$ for the BKOFF dataset

Recall the dataset Bernard & Killworth Office. Weighted graph: Ordered m = 238 edges for n = 40 nodes. The plot of  $X_3$  the number of 3-cliques:



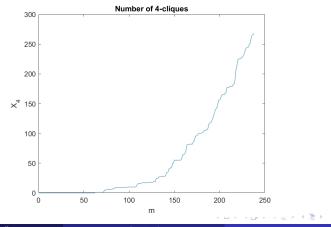
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Numerical results Plot of X<sub>4</sub> for the BKOFF dataset

Weighted graph: Ordered m = 238 edges for n = 40 nodes. The plot of  $X_4$  the number of 4-cliques:



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