Lecture 1: Graphs, Adjacency Matrices, Graph Laplacian

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Matrix Analysis

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- Directed Graphs: In a directed graph, edges are oriented. In general (i, j) ≁ (j, i).
- Weighted Graphs: Each edge has an associated weight. A weighted graph is defined by a triple (V, E, w), where w : E → ℝ is a weight function.

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Definitions Paths

Concept: A path is a sequence of edges where the right vertex of one edge coincides with the left vertex of the following edge. Example:



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A complete graph with *n* vertices has $m = \begin{pmatrix} n \\ 2 \end{pmatrix} = \frac{n(n-1)}{2}$ edges.

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Definitions Graph Attributes

Example:



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Definitions Graph Attributes

Example:



- This graph is not connected.
- It is not complete.
- It is the union of two connected graphs.

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Definitions Metric

Distance between vertices: For two vertices x, y, the distance d(x, y) is the length of the shortest path connecting x and y. If x = y then d(x, x) = 0.

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Distance between vertices: For two vertices x, y, the distance d(x, y) is the length of the shortest path connecting x and y. If x = y then d(x, x) = 0. In a connected graph the distance between any two vertices is finite. In a complete graph the distance between any two distinct vertices is 1. The converses are also true:

- **1** If $\forall x, y \in \mathcal{E}$, $d(x, y) < \infty$ then $(\mathcal{V}, \mathcal{E})$ is connected.
- **2** If $\forall x \neq y \in \mathcal{E}$, d(x, y) = 1 then $(\mathcal{V}, \mathcal{E})$ is complete.

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Definitions Metric

Graph diameter: The diameter of a graph $G = (\mathcal{V}, \mathcal{E})$ is the largest distance between two vertices of the graph:

$$D(G) = \max_{x,y \in \mathcal{V}} d(x,y)$$

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Definitions The Adjacency Matrix

For a graph $G = (\mathcal{V}, \mathcal{E})$ the adjacency matrix is the $n \times n$ matrix A defined by:

$$A_{i,j} = \left\{ egin{array}{ccc} 1 & \textit{if} & (i,j) \in \mathcal{E} \\ 0 & \textit{otherwise} \end{array}
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For directed graphs the adjacency matrix may not be symmetric. For weighted graphs $G = (\mathcal{V}, \mathcal{E}, W)$, the weight matrix W is simply given by

$$W_{i,j} = \left\{egin{array}{cc} w_{i,j} & ext{if} & (i,j) \in \mathcal{E} \ 0 & ext{otherwise} \end{array}
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Vertex Degree d(v)

For an undirected graph $G = (\mathcal{V}, \mathcal{E})$, let d(v) denote the number of edges at vertex $v \in \mathcal{V}$. The number d(v) is called the degree (or valency) of vertex v.

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$$d(v) = 2$$
, $\forall v$

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Vertex Degree Matrix D

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Graph Laplacian △

For a graph $G = (\mathcal{V}, \mathcal{E})$ the graph Laplacian is the $n \times n$ symmetric matrix Δ defined by:

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Graph Laplacian Example



Matrix Analysis

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Normalized Laplacians $\tilde{\Delta}$

Normalized Laplacian: (using pseudo-inverses)

$$\begin{split} \tilde{\Delta} &= D^{-1/2} \Delta D^{-1/2} = I - D^{-1/2} A D^{-1/2} \\ \tilde{\Delta}_{i,j} &= \begin{cases} 1 & \text{if } i = j \text{ and } d_i > 0 \\ -\frac{1}{\sqrt{d(i)d(j)}} & \text{if } (i,j) \in \mathcal{E} \\ 0 & \text{otherwise} \end{cases} \end{split}$$

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Note:

$$\Delta D^{-1} = I - AD^{-1} = L^{T} \quad ; \quad (D^{-1})_{kk} = (D^{-1/2})_{kk} = 0 \quad \text{if } d(k) = 0$$

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Image: A matrix and a matrix

Normalized Laplacians Example



Matrix Analysis

Normalized Laplacians Example



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Spectral Analysis Eigenvalues and Eigenvectors

Recall the eigenvalues of a matrix T are the zeros of the characteristic polynomial:

$$p_T(z) = det(zI - T) = 0.$$

There are exactly *n* eigenvalues (including multiplicities) for a $n \times n$ matrix *T*. The set of eigenvalues is calles its *spectrum*.

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Remark. Since $det(A_1A_2) = det(A_1)det(A_2)$ and $L = D^{-1/2}\tilde{\Delta}D^{1/2}$ it follows that $eigs(\tilde{\Delta}) = eigs(L) = eigs(L^T)$.

Spectral Analysis Rayleigh Quotient

Recall the following result:

Theorem

Assume T is a real symmetric $n \times n$ matrix. Then:

- All eigenvalues of T are real numbers.
- ② There are n eigenvectors that can be normalized to form an orthonormal basis for ℝⁿ.
- **3** The largest eigenvalue λ_{max} and the smallest eigenvalue λ_{min} satisfy

$$\lambda_{max} = \max_{x \neq 0} \frac{\langle Tx, x \rangle}{\langle x, x \rangle} \quad , \quad \lambda_{min} = \min_{x \neq 0} \frac{\langle Tx, x \rangle}{\langle x, x \rangle}$$

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Spectral Analysis Rayleigh Quotient

For two symmetric matrices T, S we say $T \leq S$ if $\langle Tx, x \rangle \leq \langle Sx, x \rangle$ for all $x \in \mathbb{R}^n$.

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 $\lambda_{\min} I \leq T \leq \lambda_{\max} I$

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In particular we say T is positive semidefinite $T \ge 0$ if $\langle Tx, x \rangle \ge 0$ for every x.

It follows that T is positive semidefinite if and only if every eigenvalue of T is positive semidefinite (i.e. non-negative).

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